

Remote Incremental Adaptation of Sensor Network Applications

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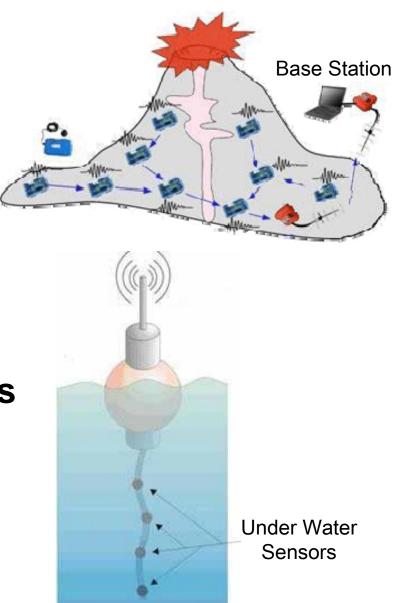


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Motivation

- WSN require uninterrupted operation indefinitely
- Customizing the system to the environment
 - Feature upgrades
 - Retasking of the System
- Post-deployment software updates are common
 - Bug Removal





State of the art

- Image replacement
- Virtual machines
- Dynamic OS's

Disadvantages:

Energy wastage, Restricted reconfigurability, Clean slate approaches

- TinyOS A framework to generate application specific OS
 - Component based architecture
 - Select app components, statically analyze and optimize
 - Supports only full binary upgrades

Goals

Energy Efficient, Fully configurable, Integrated and Transparent, Fulfilling usual Embedded System Constraints



Outline

- Motivation
- Strategy
- Architecture
 - Overview
 - Components
 - Update Procedure

Evaluation

- Transfer cost
- Update cost

Conclusion & Future Work

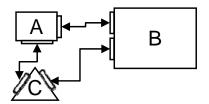


Strategy

TinyOS as a base

- Seasoned code repository
- Wide user base
- Runtime wiring

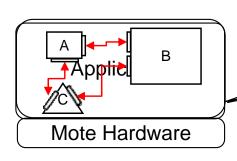
Mote Reconfiguration



Appropriate Components wired together



Compiled + Linked + Loaded



Mote reconfigured

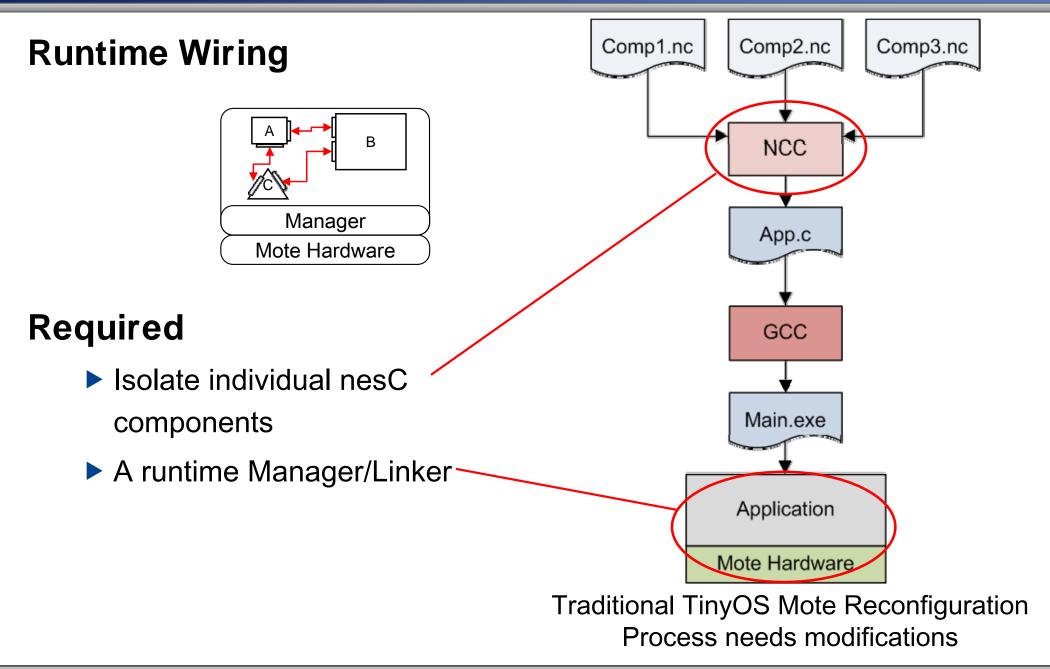
with a new

Monolithic Binary

Component structure
 Preserved

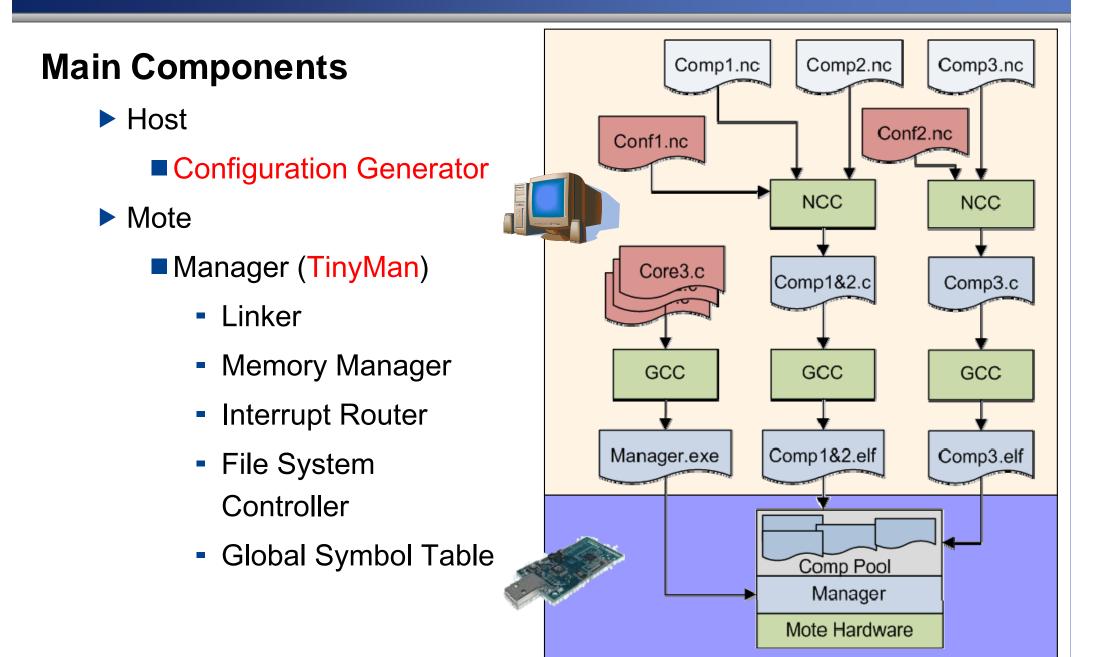


Architecture – Overview





Architecture – Overview contd.

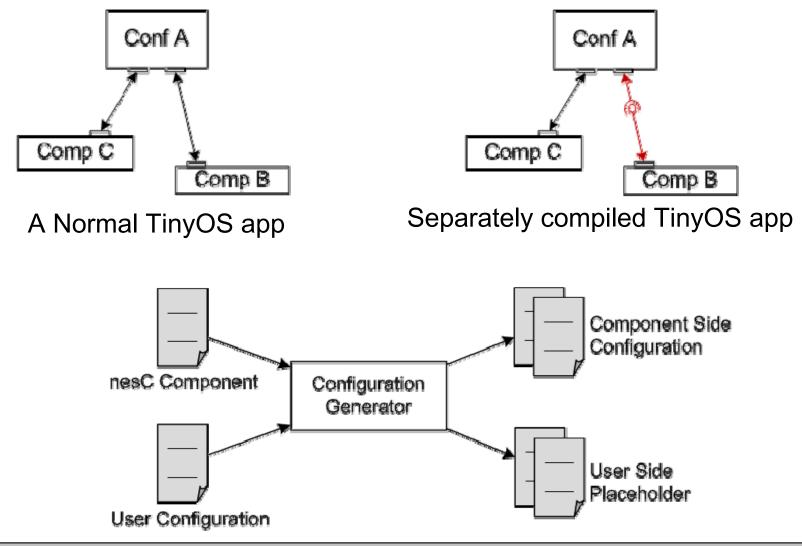




Architecture – Host

Configuration Generator

Disambiguates NCC's operation





Architecture – Mote

TinyMan

- Linker
 - Memory Manager: Manages allocation of internal flash
 - Loader: loads the linked Objects into program memory
- Interrupt Router

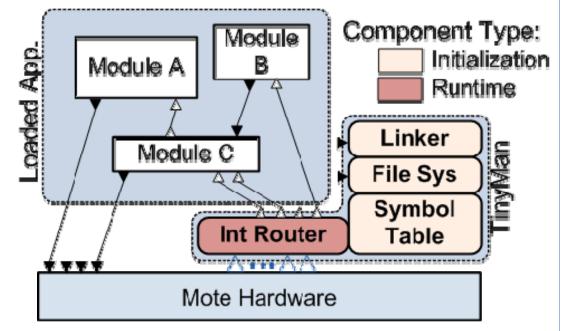
Reroutes interrupts to the functions registered as ISRs

File System

Manages storage and retrieval of ELF objects to and from non-volatile storage.

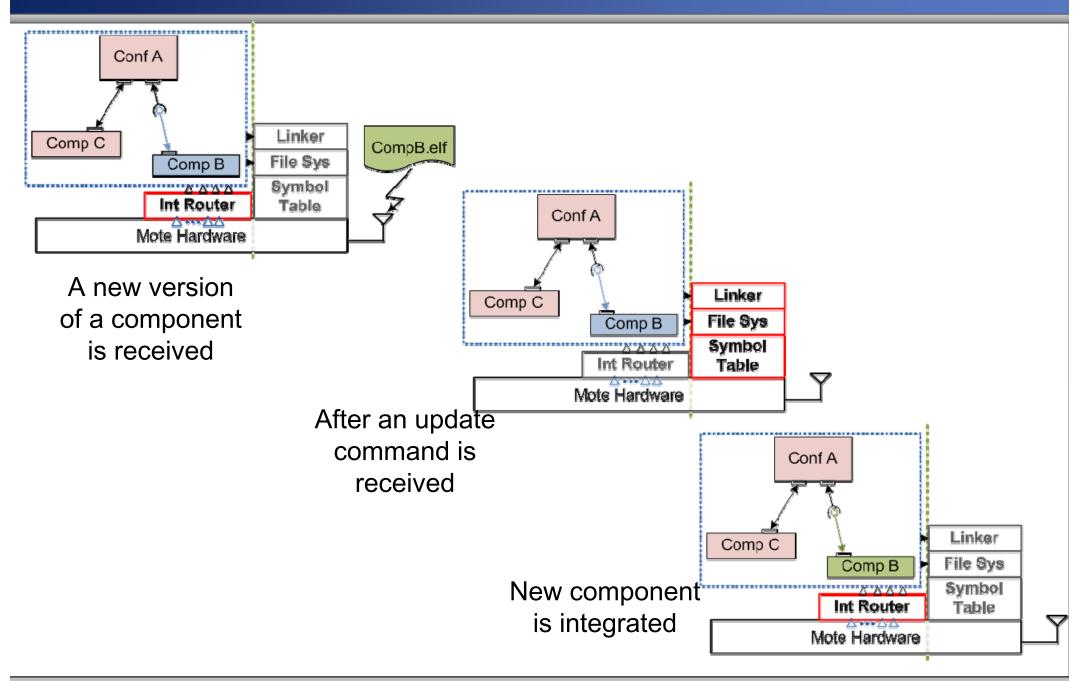
Symbol Table

Manages list of available symbols for cross linking support





Architecture – Update Procedure

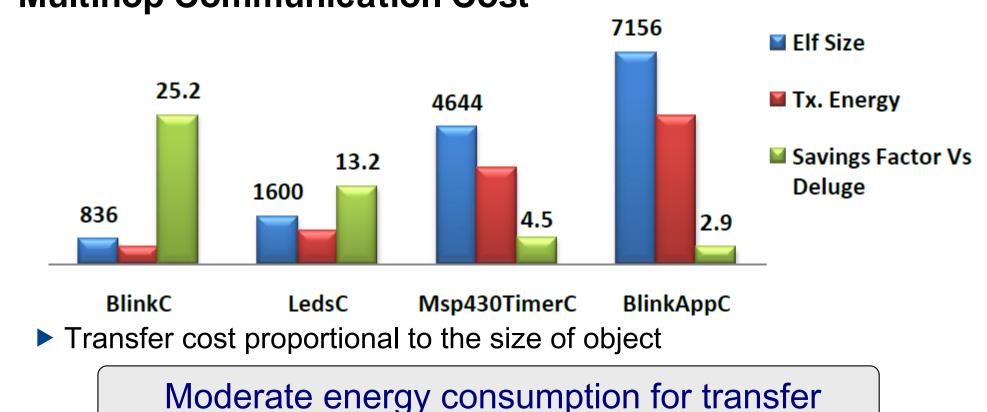




Evaluation – Update Cost

Update Cost in terms of Energy Include

- Multihop communication cost
- Cost of processing a received update



Multihop Communication Cost

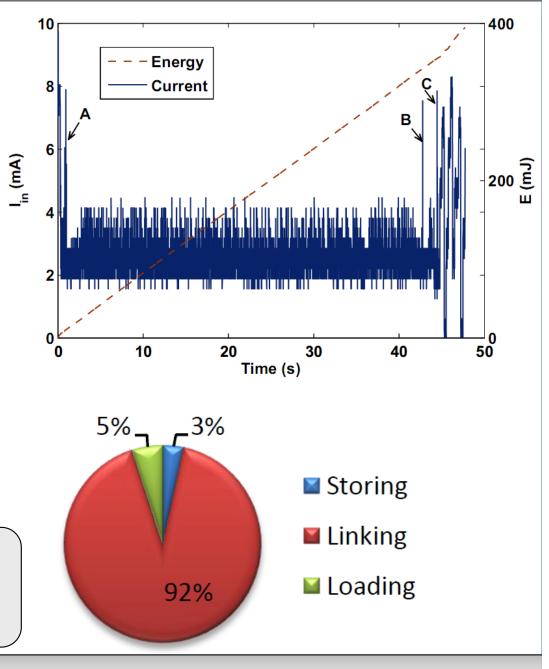


Evaluation – Update Cost contd.

Update Processing Cost

- Linking consumes most energy
 - Exponentially proportional number of string comparison operations
- Does not depend on size of update

Linking: prime candidate for optimization





Future Work & Conclusion

Conclusion

- Sensor network applications reconfigured
 - Remotely, Incrementally, With selectable granularity
- Transparent operation user friendly
- Reasonable Memory footprint
- Almost zero performance overhead

Future Work

- Automating the Configuration Generator
- ELF format optimization
 - Size Reduction due to 16 bit addressing is 8%
- Wiring compatibility checks
- Versioning scheme

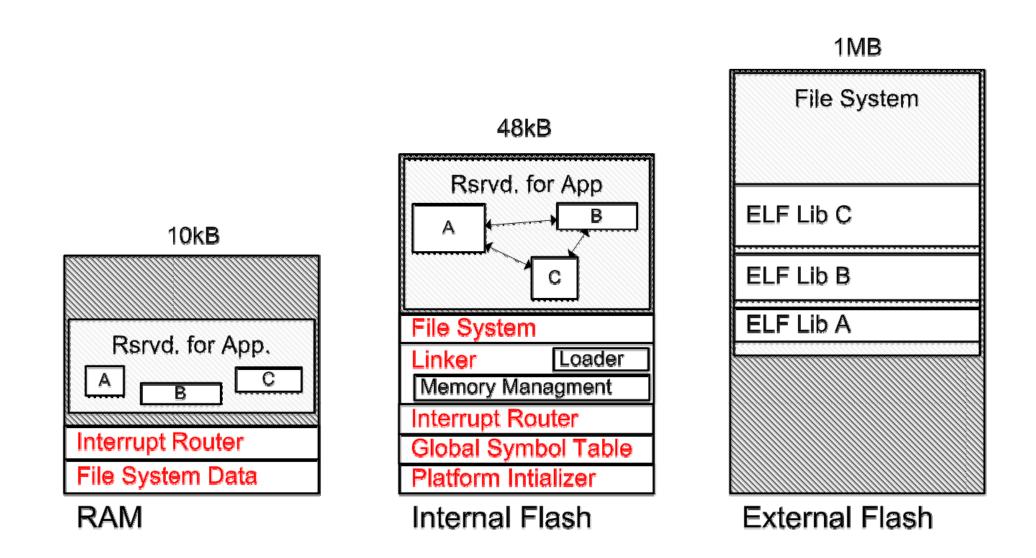




Thank you Questions ?









Related Work

Category	Solution	Update Mechanism	Pros /Cons	
Image	Хпр		+ No linking Required + Relatively transparent operation	
Replacement	Deluge	New Binary Image	 + No Execution Overhead - Large communication overhead 	
Virtual Machine	Maté	New Script	 + small communication overhead - Coarse configuration granularity - Large execution overhead 	
Dynamic Operating System	Impala	New Application	- Fixed kernel + App side API	
	Contiki	New ELF File	 Independent from TinyOS Multiple ELFs not supported 	
	SOS	New PIC Module	- PIC modules	
	FlexCup	New Module	- Not compatible with new nesC features	
	FiGaRo	New Module	- Clean slate approach, A new C API used	



Evaluation

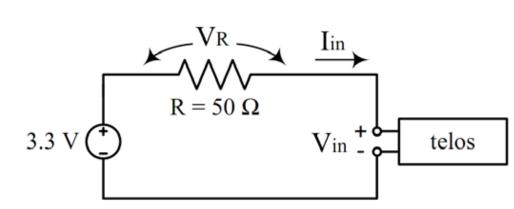
Evaluation Criteria

- Update Cost
- ELF format suitability
- Memory Footprint
- Performance Overhead

Setup

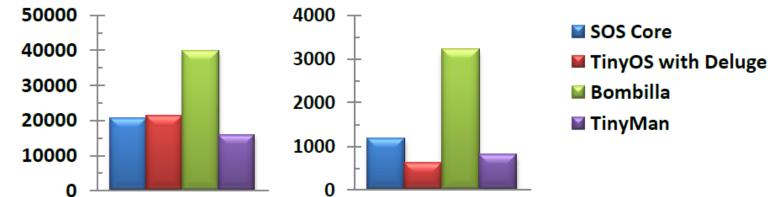
- Platform used: telos rev B
- Applications Used
 - FFT Calculator: A processor intensive application with no IO
 - Blink: An IO intensive application, Uses Timers
 - Radio ping: Application to analyze Multihop Transfer Costs





Memory Footprint - On telosB

7.7 % of RAM, 32 % of total program memory, 1.4% of total non-volatile storage



RON	I R	AM
System Name	Flash ROM (bytes)	RAM (bytes)
SOS Core	20464	1163
TinyOS with Deluge	21132	597
Bombilla Virtual Machine	39746	3196
TinyMan	15826	792

Performance Overhead

A worst case delay of 23 cycles in interrupt processing

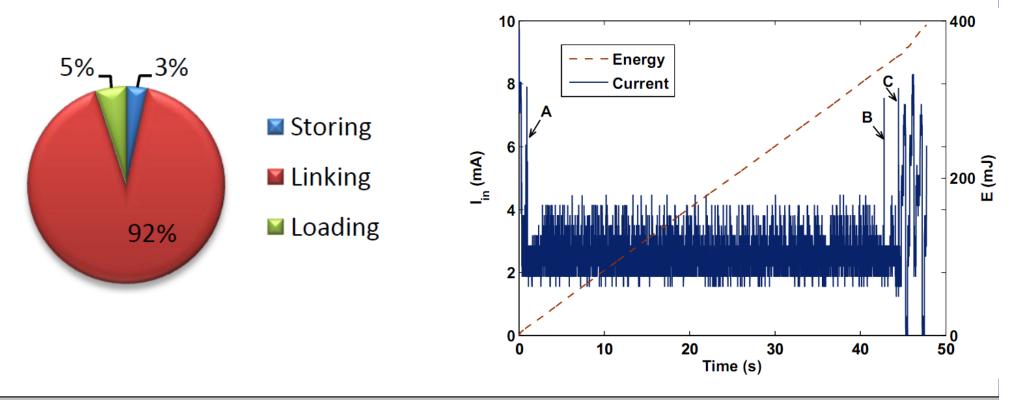


App.	Components	ELF Size	Tx. Energy (mJ)	Saving Factor (Vs. Deluge)
	BlinkC	836	58.77	25.2
	LedsC	1600	112.48	13.2
Blink	Msp430TimerC	4644	326.47	4.5
	BlinkAppC	7156	503.06	2.9
	Tota	1	1000.78	
			7156	Elf Size
25	.2	4644		🖬 Tx. Energy
836	13.2 1600	4.5		Savings Factor* 2.9
BlinkC	LedsC	Msp430Time	erC BlinkAp	рС



Update Processing Cost

Application	Stong	Size	Time	Energy Consumed
Application	Steps	(B)	(s)	(mJ)
Blink	Storing	14236	0.9	11.4
	Linking	14236	43.2	305.1
	Loading	4756	1.5	16.7
	Total		45.6	333.2





Detailed Memory Footprint

Component		Flash ROM (bytes)	RAM (bytes)
File System		5630	228
Platform Initializer		454 0	
	Linker Main	368	4
Linker	Memory Management	344	84
	Loader	2460	24
Interrupt Router		1092	30
Global Symbol Table		1028	74
Node Runtime		390	0
	UART	102	259
	Flash Write	5630 454 368 344 2460 1092 1028 390	4
Hardware Drivers	SPI	54	0
	Ex Flash	726	0
Library Reference		3020	85
Total		15826	792



ELF Sections Detail

